The Heisenberg–Weil Representation and Fast Wireless Communication

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Madison

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May 12, 2012 1 / 28

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(0) Motivation - GPS

• GPS



CLIENT WANT: Coordinates of satellite and time delay (enables to calculate distance to a satellite)?

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May 12, 2012 3 / 28

• S, $R \in \mathcal{H} = \mathbb{C}(\mathbb{Z}_N)$ – Hilbert space of digital signals, $N \gg 1000$.

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S, R ∈ H = C(Z_N) - Hilbert space of digital signals, N ≫ 1000. S, R : {0, ..., N − 1} → C.

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- S, $R \in \mathcal{H} = \mathbb{C}(\mathbb{Z}_N)$ Hilbert space of digital signals, $N \gg 1000$.
 - $S, R: \{0, \dots, N-1\} \rightarrow \mathbb{C}.$
 - Satellite transmits $b \cdot S$, b =coordinates.

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Fact (GPS)

Client receives

$$R[n] = b \cdot \sum_{k=1}^{m} \alpha_k \cdot e^{\frac{2\pi i}{N} \omega_k \cdot n} \cdot S[n + \tau_k] + \mathcal{W}[n], \quad n \in \mathbb{Z}_N,$$

m = # paths, $\alpha_k \in \mathbb{C}$ intensity, $\sum_{k=1}^m |\alpha_k|^2 \leq 1$, $\omega_k \in \mathbb{Z}_N$ Doppler, $\tau_k \in \mathbb{Z}_N$ delay, along path k, $\mathcal{W} \in \mathcal{H}$ random noise.

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Problem

Design $S \in \mathcal{H}$, and effective method to extract (b, τ) , $\tau = \min{\{\tau_k\}}$, using R and S.

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• Simpler scenario

$$R[n] = e^{\frac{2\pi i}{N}\omega_0 \cdot n} \cdot S[n+\tau_0] + \mathcal{W}[n].$$

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$$R[n] = e^{\frac{2\pi i}{N}\omega_0 \cdot n} \cdot S[n+\tau_0] + \mathcal{W}[n].$$

Problem (Time-Frequency Shift)

Design $S \in \mathcal{H}$, and method of extracting (τ_0, ω_0) from S and R.

(I) Solution - MATCHED FILTER

Definition

Matched filter

$$\mathcal{M}(R,S): \underbrace{\mathbb{Z}_{N} \times \mathbb{Z}_{N}}_{\mathsf{Time-Frequency}} \to \mathbb{C},$$
$$\mathcal{M}(R,S)[\tau,\omega] = \left\langle R[n] , e^{\frac{2\pi i}{N}\omega \cdot n} \cdot S[n+\tau] \right\rangle.$$

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Identity

$$\mathcal{M}(R,S)[\tau,\omega] = \mathcal{M}(S,S)[\tau-\tau_0,\omega-\omega_0] + O(\frac{NSR}{\sqrt{N}}).$$

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Identity

$$\mathcal{M}(R,S)[\tau,\omega] = \mathcal{M}(S,S)[\tau-\tau_0,\omega-\omega_0] + O(\frac{NSR}{\sqrt{N}}).$$

• Question: What S to use for extracting (au_0, ω_0) from $\mathcal{M}(R, S)$?

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Solution - MATCHED FILTER

• Typical solution: S = pseudo-random.

Example



 $|\mathcal{M}(\textit{R},\textit{S})|$, S=pseudo-random, $(au_{0},\omega_{0})=($ 50, 50).

• Complexity of Pixel-by-Pixel Algorithm



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• Complexity of Pixel-by-Pixel Algorithm



Problem

Faster algorithm.

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May 12, 2012 8 / 28

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• Suppose $L \subset \mathbb{Z}_N \times \mathbb{Z}_N$ line



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• Compute $\mathcal{M}(R, S)$ on L in $O(N \cdot \log(N))$ operations!

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- Example:

 $\mathcal{M}(\textit{R},\textit{S})[\tau,0] = \langle \textit{R}[\textit{n}],\textit{S}[\textit{n}+\tau] \rangle = (\textit{S}*\textit{R})[\tau] ~~\text{-}~\text{Fast by FFT}.$

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• Compute entire $\mathcal{M}(R, S)$ in $O(N^2 \cdot \log(N))$ operations.

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- Compute entire $\mathcal{M}(R, S)$ in $O(N^2 \cdot \log(N))$ operations.
- **Question:** Can you design S and method to make almost linear number of operations?

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(II) Flag Algorithm - IDEA

 Suppose for a line L ⊂ Z_N × Z_N we construct a signal S_L ∈ H with *M*(R, S_L) of the form



 $|\mathcal{M}(\textit{R},\textit{S}_L)|$, $(au_0,\omega_0)=(50,50)$

Then we have algorithm of complexity $O(N \cdot \log(N))$!!

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May 12, 2012 10 / 28

Flag Algorithm - WAVEFORM DESIGN PROBLEM

Problem (Waveform Design)

For every $L \subset \mathbb{Z}_N \times \mathbb{Z}_N$ construct $S_L \in \mathcal{H}$ with

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For every $L \subset \mathbb{Z}_N \times \mathbb{Z}_N$ construct $S_L \in \mathcal{H}$ with

1 Flag. Matched filter $\mathcal{M}(R, S_L)$ of the form



 $|\mathcal{M}(R, S_L)|$

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May 12, 2012 11 / 28

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For every $L \subset \mathbb{Z}_N \times \mathbb{Z}_N$ construct $S_L \in \mathcal{H}$ with

1 Flag. Matched filter $\mathcal{M}(R, S_L)$ of the form



 $|\mathcal{M}(R, S_L)|$

2 Almost orthogonality. For $L \neq M$ the cross-correlations $|\mathcal{M}(S_L, S_M)[\tau, \omega]| = O(\frac{1}{\sqrt{N}}).$

(III) Waveform Design - EXAMPLE

• Consider waveform
$$f[n] = \frac{1}{\sqrt{N}} e^{\frac{2\pi i}{N}n}$$

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(III) Waveform Design - EXAMPLE

• Consider waveform
$$f[n] = \frac{1}{\sqrt{N}} e^{\frac{2\pi i}{N}n}$$
.
• $f[n+\tau] = e^{\frac{2\pi i}{N}\tau} \cdot f[n]$, so
 $\mathcal{M}(f,f)[\tau,\omega] = \langle f[n], e^{\frac{2\pi i}{N}\omega \cdot n} \cdot f[n+\tau] \rangle$
 $= \begin{cases} |\cdot| = 1, \text{ if } \omega = 0; \\ 0, & \text{otherwise.} \end{cases}$

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 $= \begin{cases} |\cdot| = 1, \text{ if } \omega = 0; \\ 0, & \text{otherwise.} \end{cases}$
• Take $S = \underbrace{f}_{e \times p} + \underbrace{\varphi}_{pseudo-random}$, then
 $|\mathcal{M}(R,S)[\tau,\omega]| \approx \begin{cases} 2, & \text{if } (\tau,\omega) = (\tau_0,\omega_0); \\ 1, & \text{on the line } \omega = \omega_0; \\ O(\frac{1}{\sqrt{N}}), & \text{otherwise.} \end{cases}$

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FLAG - Numerics



May 12, 2012 13 / 28

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 \bullet Question: How to generalize the "good" orthonormal basis of ${\cal H}$

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• Answer: Consider the Heisenberg operators on $\mathcal{H} = \mathbb{C}(\mathbb{Z}_N)$

$$\begin{cases} \pi: \mathbb{Z}_N \times \mathbb{Z}_N \to U(\mathcal{H}), \\ \pi(\tau, \omega) f[n] = e^{\frac{2\pi i}{N}\omega \cdot n} \cdot f[n+\tau]. \end{cases}$$

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Property:

$$\pi(\tau,\omega)\circ\pi(\tau',\omega')=e^{\frac{2\pi i}{N}(\tau\omega'-\omega\tau')}\cdot\pi(\tau',\omega')\circ\pi(\tau,\omega).$$

Restrict π to the line L = {(τ, 0); τ ∈ Z_N}, obtain commutative collection of operators

time shift $\overbrace{\pi(\tau,0)}^{\sim}:\mathcal{H}\to\mathcal{H}, \ \tau\in\mathbb{Z}_N.$

Restrict π to the line L = {(τ, 0); τ ∈ Z_N}, obtain commutative collection of operators

$$\overbrace{\pi(\tau,0)}^{\text{time shift}}:\mathcal{H}\to\mathcal{H}, \ \tau\in\mathbb{Z}_N.$$

Theorem (Linear Algebra – Simultaneous Diagonalization)

There exists a basis for $\mathcal H$ of common eigenfunctions

$$\mathcal{B}_L = \{ f_{\psi}; \ \pi(\tau, 0)[f_{\psi}] = \overbrace{\psi(\tau)}^{e.v.} \cdot f_{\psi}, \ \text{for all } \tau \in \mathbb{Z}_N \}.$$

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• Of course $\mathcal{B}_L = \mathcal{B}_L$.

• Mechanism to generate 'good waveforms':

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May 12, 2012 16 / 28
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Theorem (Support, [Calderbank–Howard–Moran, Howe])

We have

$$\left|\mathcal{M}(f_{\psi},f_{\psi})[\tau,\omega]\right| = \begin{cases} 1, \ (\tau,\omega) \in L; \\ 0, \ (\tau,\omega) \notin L. \end{cases}$$

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HEISENBERG (LINES) SYSTEM - Numerics



 $\left|\mathcal{M}[f_{\psi},f_{\psi}]\right|, f_{\psi} \in \mathcal{B}_{L}, L = \{(\tau,\tau)\}$

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May 12, 2012 17 / 28

HEISENBERG (LINES) SYSTEM – Numerics



• What is next?

HEISENBERG (LINES) SYSTEM - Numerics



$$\left|\mathcal{M}[f_{\psi}, f_{\psi}]\right|, f_{\psi} \in \mathcal{B}_L, L = \{(\tau, \tau)\}$$

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- What is next?
- Just add any pseudo-random waveform.

• Discrete Fourier transform (DFT) is defined by system

$$\Sigma_W: \quad DFT \circ \pi \begin{pmatrix} \tau \\ \omega \end{pmatrix} = \pi(\underbrace{\begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}}_{W} \begin{pmatrix} \tau \\ \omega \end{pmatrix}) \circ DFT, \quad \tau, \omega \in \mathbb{Z}_N.$$

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Special linear group

$$W \in SL_2(\mathbb{Z}_N) = \{g = \begin{pmatrix} a & b \\ c & d \end{pmatrix}; a, b, c, d \in \mathbb{Z}_N, \det(g) = 1\}.$$

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Problem (André Weil)

For each $g \in SL_2(\mathbb{Z}_N)$ find operator $\rho(g)$ on $\mathcal{H} = \mathbb{C}(\mathbb{Z}_N)$, which solves the system of N^2 linear conditions

$$\Sigma_{g}: \qquad \underbrace{\rho(g)}_{?} \circ \pi\begin{pmatrix} \tau\\ \omega \end{pmatrix} = \pi(g \cdot \begin{pmatrix} \tau\\ \omega \end{pmatrix}) \circ \underbrace{\rho(g)}_{?}, \qquad \tau, \omega \in \mathbb{Z}_{N}.$$

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Theorem (Stone-von Neumann-Schur-Weil)

 $\exists ! \ \textit{collection} \ \{\rho(g) \in Sol(\Sigma_g); \ g \in \textit{SL}_2(\mathbb{Z}_N) \}$ such that

$$\begin{cases} \rho: SL_2(\mathbb{Z}_N) \to U(\mathcal{H}), \\ \rho(g \cdot h) = \rho(g) \circ \rho(h). \end{cases}$$
 — Weil representation

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• Mechanism to generate 'good waveforms':

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- Mechanism to generate 'good waveforms':
 - Take maximal commutative subgroup T ⊂ SL₂(Z_N), i.e., gh = hg for every g, h ∈ T.

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 - Get a <u>commutative</u> collection of operators

$$\rho(g): \mathcal{H} \to \mathcal{H}, g \in T.$$

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 $\bullet\,$ Get a basis for ${\cal H}$ of common eigenfunctions

$$\mathcal{B}_{\mathcal{T}} = \{ \varphi_{\chi}; \ \rho(g)[\ \varphi_{\chi}] = \overbrace{\chi(g)}^{\text{e.v.}} \cdot \varphi_{\chi}, \ \text{ for all } g \in \mathcal{T} \}.$$

Theorem (Pseudo-Randomness [G–Hadani–Sochen])

For
$$\varphi_{\chi} \in \mathcal{B}_{T}$$
 we have $\left| \mathcal{M}(\varphi_{\chi}, \varphi_{\chi})[\tau, \omega] \right| = \begin{cases} 1, & (\tau, \omega) = (0, 0); \\ \leq \frac{2}{\sqrt{N}}, & (\tau, \omega) \neq (0, 0). \end{cases}$



Theorem ([Fish–G–Hadani–Sayeed–Schwartz])

Take
$$S_L = \underbrace{f_L}_{\in \mathcal{B}_L} + \underbrace{\varphi_{\chi}}_{\in \mathcal{B}_T}$$
. Then

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Theorem ([Fish–G–Hadani–Sayeed–Schwartz])

Take $S_L = \underbrace{f_L}_{\in \mathcal{B}_L} + \underbrace{\varphi_{\chi}}_{\in \mathcal{B}_T}$. Then

• Flag. We have
$$|\mathcal{M}(S_L, S_L)[\tau, \omega]| = \begin{cases} \approx 2, \text{ if } (\tau, \omega) = (0, 0); \\ \approx 1, \text{ if } (\tau, \omega) \in L \setminus (0, 0); \\ \leq \frac{7}{\sqrt{N}}, \text{ otherwise.} \end{cases}$$

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2 Almost orthogonality. For $L \neq M$ we have $|\mathcal{M}(S_L, S_M)[\tau, \omega]| \leq \frac{7}{\sqrt{N}}$, for every $(\tau, \omega) \in \mathbb{Z}_N \times \mathbb{Z}_N$.

Theorem ([Fish–G–Hadani–Sayeed–Schwartz])

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(IV) APPLICATIONS

- (A) Channel Estimation
 - Transmit flag waveform $S = S_L$



Image: Image:

(IV) APPLICATIONS

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Receive

$$R[n] = \sum_{k=1}^{m} \alpha_k \cdot \underbrace{e^{\frac{2\pi i}{N}\omega_k \cdot n}}_{\text{Doppler}} \cdot S_L[n + \underbrace{\tau_k}_{\text{delay}}] + \mathcal{W}[n],$$

 $\mathcal{W} = \text{noise}, \ \tau_k = \text{delay along path } k, \ \omega_k = \text{Doppler along path } k, \ \alpha_k = \text{intensity coefficient}, \ |\alpha_1|^2 + ... + \ |\alpha_m|^2 \leq 1.$

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$$R[n] = \sum_{k=1}^{m} \alpha_k \cdot \underbrace{e^{\frac{2\pi i}{N}\omega_k \cdot n}}_{\text{Doppler}} \cdot S_L[n + \underbrace{\tau_k}_{\text{delay}}] + \mathcal{W}[n],$$

 $\mathcal{W} =$ noise, $\tau_k =$ delay along path k, $\omega_k =$ Doppler along path k, $\alpha_k =$ intensity coefficient, $|\alpha_1|^2 + ... + |\alpha_m|^2 \leq 1$. • Goal Extract $(\alpha_k, \tau_k, \omega_k)$'s, using R and S_L .

Application - CHANNEL ESTIMATION

• We have
$$|\mathcal{M}(R, S_L)[\tau, \omega]| \approx$$

$$\begin{cases}
2 \cdot \alpha_k, \text{ if } (\tau, \omega) = (\tau_k, \omega_k); \\
1 \cdot \alpha_k, \text{ if } (\tau, \omega) \in L + (\tau_k, \omega_k) \smallsetminus (\tau_k, \omega_k); \\
\leq O(\frac{m}{\sqrt{N}}), \text{ otherwise.} \end{cases}$$



 $|\mathcal{M}(R, S_L)|, L = \{(\tau, 0)\}, (\alpha_k, \tau_k, \omega_k) = (\frac{1}{\sqrt{3}}, 50k, 50k), k = 1, 2, 3.$

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1 \cdot \alpha_k, \text{ if } (\tau, \omega) \in L + (\tau_k, \omega_k) \smallsetminus (\tau_k, \omega_k); \\
\leq O(\frac{m}{\sqrt{N}}), \text{ otherwise.} \end{cases}$$



 $|\mathcal{M}(R, S_L)|, L = \{(\tau, 0)\}, (\alpha_k, \tau_k, \omega_k) = (\frac{1}{\sqrt{3}}, 50k, 50k), k = 1, 2, 3.$

• Flag method computes channel parameters in $O(m \cdot N \log(N))$.

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Fast Wireless Communication

May 12, 2012 23 / 28

Application - MOBILE COMMUNICATION

• (B) Mobile Communication: CDMA, GSM...



Application - MOBILE COMMUNICATION

• (B) Mobile Communication: CDMA, GSM...



• User transmits message

 $b^{+1 \text{ or } -1} \cdot S_l$

using his private flag waveform S_L .

Application - MOBILE COMMUNICATION

• Antenna receives:

$$R[n] = b \cdot \sum_{k=1}^{m} \alpha_k \cdot e^{\frac{2\pi i}{N}\omega_k \cdot n} \cdot S_L[n+\tau_k] + \mathcal{W}[n],$$

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• Assumption (slow varying channel): Antenna knows, using our $O(m \cdot N \log(N))$ channel estimation, the $(\alpha_k, \omega_k, \tau_k)$'s.

• Antenna receives:

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- Goal: Extract b using R and S_L .

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- Goal: Extract b using R and S_L .
- Method:

$$\left\langle R[n], \sum_{k=1}^{m} \alpha_k \cdot e^{\frac{2\pi i}{N}\omega_k \cdot n} \cdot S_L[n+\tau_k] \right\rangle \approx \left(\sum_{k=1}^{m} |\alpha_k|^2 \right) \cdot b.$$

• (C) GPS



CLIENT WANT: Coordinates of satellite and time delay (enables to calculate distance to a satellite)?

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May 12, 2012 26 / 28

•
$$S_L$$
, $R \in \mathcal{H} = \mathbb{C}(\mathbb{Z}_N)$, $N >> 1000$.

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Fact

GPS System

$$R[n] = b \cdot \sum_{k=1}^{m} \alpha_k \cdot e^{\frac{2\pi i}{N}\omega_k \cdot n} \cdot S_L[n+\tau_k] + \mathcal{W}[n],$$

 $\tau_k = delay, \ \omega_k = Doppler, \ b = coordinates \ of \ satellite, \ \alpha_k = intensity \ coefficient, \ |\alpha_1|^2 + ... + |\alpha_m|^2 \leq 1.$

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Problem

Extract
$$(b, \tau)$$
, $\tau = \min{\{\tau_k \ s\}}$ using R and S_L .

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GPS System

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Problem

Extract
$$(b, \tau)$$
, $\tau = \min{\{\tau_k \ s\}}$ using R and S_L .

Solution

Flag method solves in $O(m \cdot N \log(N))$.

Shamgar Gurevich (Madison)
THANK YOU



Sasha

Ronny



Akbar

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Shamgar Gurevich (Madison)

æ May 12, 2012

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